

Goal:

Simplify information processing tasks by reducing permutation invariant systems to simple de Finetti systems.

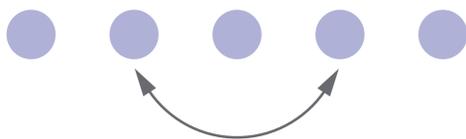
Theorem:

There exists a de Finetti system, $\tau_{A|X}$, such that for every permutation invariant system $P_{A|X}$

$$\forall a, x \quad P_{A|X}(a|x) \leq (n+1)^{m(l-1)} \tau_{A|X}(a|x).$$

Permutation invariant systems:

We can permute the subsystems and all will stay the same



de Finetti systems:

Systems with a **simple** structure - a convex combination of i.i.d. systems

Operational definition of a system:

- Conditional probability distributions $P_{A|X}$
- X - measurements
- A - outcomes
- Describes a larger set of systems than quantum systems



de Finetti reductions for conditional probability distributions

- Reduction from permutation invariant conditional probability distributions to a de Finetti system
- Independent of the **dimension** of the underlying space
- Depends on the **number of measurements and outcomes**
- **Useful for device independent tasks**

- $\tau_{A|X}$ - de Finetti system
- n - # of subsystems
- m - # of measurements of each subsystem
- l - # of outcomes of each subsystem



Applications

Instead of analyzing lower bounds for every possible system — analyze it only for the simple de Finetti system and “pay” for it with a polynomial factor

Lemma:

Consider a permutation invariant test which interacts with a system $P_{A|X}$ and outputs “success” or “fail” with some probabilities. Then for every system $P_{A|X}$

$$\Pr_{\text{fail}}(P_{A|X}) \leq (n+1)^{m(l-1)} \Pr_{\text{fail}}(\tau_{A|X}).$$

More on the arXiv:

- Similar lower bounds on the diamond norm — useful for cryptography
- Better bounds in the presence of symmetries
- **Example application** — *simplifying the analysis of CHSH based protocols*

