

# Practical Parallel Self-testing of Bell States via Magic Rectangles\*

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## Abstract

Self-testing is a method to verify that one has a particular quantum state from purely classical statistics. For applications such as device-independent delegated verifiable quantum computation, it is crucial that one tests multiple Bell states in parallel while keeping the quantum requirements of one side to a minimum. We use  $3 \times n$  magic rectangle games to obtain a self-test for  $n$  Bell states where one side need only make single-qubit Pauli measurements. It consumes little randomness, is robust, and requires only perfect correlations. To achieve this, we introduce a one-side-local quantum strategy for the magic square game that wins with certainty, generalise this to the family of  $3 \times n$  magic rectangle games, and supplement these games with extra check rounds.

## Magic square game

The *magic square* game is a nonlocal game played on a  $3 \times 3$  grid [1].

- Alice and Bob are assigned (uniformly at random) a row and column.

Players must fill their row/column with  $\pm 1$  according to certain rules:

1. The product of Alice's row must be **positive**.
2. The product of Bob's column must be **negative**.



Figure 1. Example for Alice (left) and Bob (middle). The players win (right).

**Win condition:** Values entered into the *shared cell* coincide.

## One-side-local strategy

Optimal **classical** and **quantum** win probabilities  $8/9$  and  $1$ .

- Standard strategy has  $|\Phi^+\rangle_{AB}^{\otimes 2}$  shared between Alice and Bob.
- Requires two-qubit *entangled* measurements upon some inputs.
- Alice needs only **single-qubit** measurements if  $|\Phi^+\rangle_{AB}^{\otimes 3}$  shared.

|           |          |           |          |          |          |
|-----------|----------|-----------|----------|----------|----------|
| $X_1$     | $X_1X_2$ | $X_2$     | $X_2X_3$ | $X_1X_3$ | $X_1X_2$ |
| $-X_1Z_2$ | $Y_1Y_2$ | $-Z_1X_2$ | $Y_2Y_3$ | $Y_1Y_3$ | $Y_1Y_2$ |
| $Z_2$     | $Z_1Z_2$ | $Z_1$     | $Z_2Z_3$ | $Z_1Z_3$ | $Z_1Z_2$ |

Figure 2. The *standard* strategy (left) and *one-side-local* strategy (right).

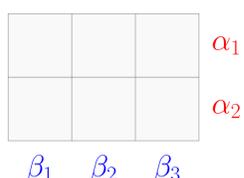
The **one-side-local** strategy generalises to  $3 \times n$  *magic rectangle* games (for  $n = 3 \pmod{4}$ ) and has a similar structure.

## Generalisation: Magic rectangle games

*Magic rectangle* games [2] are played on an  $m \times n$  grid.

The rules are generalised accordingly:

1. The product of Alice's  $i$ th row must be  $\alpha_i$ .
2. The product of Bob's  $j$ th row must be  $\beta_j$ .



To avoid deterministic winning strategies, we also require

$$\alpha_1 \dots \alpha_m \cdot \beta_1 \dots \beta_n = -1.$$

The self-test uses 3 rows,  $3 \pmod{4}$  columns,  $\alpha_i = +1$ , and  $\beta_j = -1$ .

## Self-testing protocol: Three Bell states

Let  $n = 3$  be the number of Bell states to be tested. In each round, a verifier chooses  $c \in \{0, 1\}$  and  $y \in \{1, \dots, n\}$ . The verifier sends Bob  $(c, y)$  and, depending on  $c$ , runs one of the following subprotocols:

0. **Magic game.** Send Alice  $x \in \{1, 2, 3\}$ . Alice and Bob answer with  $a_1, \dots, a_n$  and  $b_1, b_2, b_3$  in  $\{+1, -1\}$  satisfying  $b_1b_2b_3 = -1$ . Accept if and only if  $\prod_{k \neq y} a_k = b_x$ .
1. **Local check.** Send Alice  $x \in \{1, 3\}$ . Alice and Bob answer with  $a_1, \dots, a_n$  and  $b_1, \dots, b_n$  in  $\{+1, -1\}$ . If  $x = 1$ , accept if and only if  $a_y = b_y$ . If  $x = 3$ , accept if and only if  $a_j = b_j$  for all  $j \neq y$ .

## Self-testing protocol: Many Bell states

Let  $n = 3 \pmod{4}$ . The verifier chooses  $c \in \{0, 1, 2\}$  and performs the previous protocol with an additional subprotocol if  $c = 2$  is chosen:

2. **Pair check.** Send Alice  $x \in \{1, 3\}$ . Alice answers with  $a_1, \dots, a_n$ . Bob answers with  $n - 1$  bits  $b_{y-k, y+k}$  and  $b'_{y-k, y+k}$  in  $\{+1, -1\}$  (with addition taken modulo  $n$ ) for all  $k \in \{1, \dots, \frac{n-1}{2}\}$ . If  $x = 1$ , accept if and only if  $a_i a_j = b_{i,j}$  for all  $i, j$ . If  $x = 3$ , accept if and only if  $a_i a_j = b'_{i,j}$  for all  $i, j$ .

## Robustness and completeness

If a strategy is accepted with probability at least  $1 - \epsilon$ , the protocol self-tests the state  $|\Phi^+\rangle_{AB}^{\otimes n}$  with robustness  $O(n^{5/2} \sqrt{\epsilon})$ .

- Subprotocol **magic game** ensures a perfect  $3 \times n$  strategy is used.
- **Local check** rules out *entangled* measurements for Alice.
- **Pair check** rules out deterministic extensions to single-qubit strategies using smaller  $3 \times n'$  magic rectangles.

There exist strategies (based on *one-side-local* magic game strategies) that are accepted with certainty (use only perfect correlations).

- In the *honest* case, Alice needs only single-qubit Pauli measurements, while Bob requires two-qubit, entangled measurements.

## Comparison

The protocol simultaneously achieves several properties desirable in the client/server setting.

| Protocol                     | Local        | Perf. corr. | Err. tol. $\epsilon(n, \delta)$ | Input size |                  |
|------------------------------|--------------|-------------|---------------------------------|------------|------------------|
|                              |              |             |                                 | Alice      | Bob              |
| <b>This protocol</b>         | Alice        | Yes         | $O(n^{-5}\delta^2)$             | $O(1)$     | $O(\log n)$      |
| Šupić et al. (2021)          | As base test |             | N/A                             |            | $O(1)$           |
| Chao et al. (2018)           | Yes          | No          | $O(n^{-5}\delta^2)$             |            | $O(\log n)$      |
| Natarajan and Vidick (2018)  | No           | Yes         | $O(\delta^c)$                   |            | $O(\log n)$      |
| Natarajan and Vidick (2017)  | As CHSH/MS   |             | $O(\delta^{16})$                |            | $O(n)$           |
| Coladangelo (2017): MS       | No           | Yes         | $O(n^{-3}\delta^2)$             |            | $O(n)$           |
| Coladangelo (2017): CHSH     | Yes          | No          | $O(n^{-3}\delta^2)$             |            | $O(n)$           |
| Coudron and Natarajan (2016) | No           | Yes         | $O(n^{-4}\delta^4)$             |            | $O(n)$           |
| McKague (2016)               | Yes          | No          | $O(n^{-8}\delta^8)$             |            | $O(\log \log n)$ |

Sample comparisons with other protocols, including some based on the magic square (MS) game, are shown above.

## References

- [1] Padmanabhan K. Aravind. Quantum mysteries revisited again. *Am. J. Phys.*, 72(10):1303–1307, 9 2004.
- [2] Sean A. Adamson and Petros Wallden. Quantum magic rectangles: Characterization and application to certified randomness expansion. *Phys. Rev. Research*, 2(4):043317, 12 2020.