# Improved analytical bounds on delivery times of long-distance entanglement

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The fundamental distance limit for quantum key distribution due to photon loss can be overcome by intermediate nodes called quantum repeaters. We provide analytical bounds on the mean and quantiles of the entanglement delivery time for a very general class of repeater schemes, which significantly improve upon existing work. Our bounds enable the analytical assessment of repeater in the presence of time-dependent noise, such as imperfect memories, and are useful for the design and analysis of network sizes beyond the reach of numerics.



Results: analytical bounds on the

- Example quantum repeater chain scheme on 3 nodes:
- Realizations of quantum repeaters will suffer from time-dependent noise, such as memory decoherence. Specifically: many times, an entangled pair is generated which needs to wait for another pair, and decoheres during this waiting:



Span of entanglement

• To characterize time-dependent noise, we need to know:

**Problem statement:** Given a repeater chain scheme, find the probability distribution of the time until entanglement delivery by the repeaters' end nodes

## entanglement delivery time

(1) Bounds on mean delivery time. For two input entangled pairs:



(2) Exponentially-fast decaying bounds on Pr(it takes longer than time t to deliver entanglement)



**Tool**: reliability theory. A random variable T is **new-better-than-used** (NBU) if:

- Existing literature: often even mean time not known!
- We consider hierarchical repeater chain schemes (based on the BDCZ) scheme [1]), which are:
  - composed of probabilistic components (GENeration of fresh entanglement, DISTillation, SWAPping) with Pr(success)≥const
  - No two components wait for the same entangled pair before proceeding, i.e. the entanglement dependency graph is a tree

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#### $\forall x, y \ge 0: \Pr(T > x + y) \le \Pr(T > x) \cdot \Pr(T > y)$

#### Formally, we show:

**Proposition**: let T<sub>output</sub> be the time until success of an entanglement swap with success probability p, applied to two states which are produced with iid times  $T_{input}$  (also a random variable. If  $T_{input}$  is continuous and NBU, then:

1. 
$$T_{output}$$
 is NBU  
2.  $mean(T_{output}) \le \frac{3T_{input}}{2p}$   
3.  $Pr(T_{output} > t) \le exp\left(p - \frac{2pt}{3mean(T_{input})}\right)$  (and also a lower bound)

And we have similar statements for  $\geq 2$  input states.

### Application to the famous nested BDCZ scheme



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 $\times 10^{5}$ 

1.0







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Alice







BDCZ scheme: H.-J. Briegel, W. Du'r, J. I. Cirac, and P. Zoller, Phys. Rev. Lett. 81, 5932–5935 (1998)