

else anticommute

One-time memory from isolated Majorana islands Sourav Kundu[†], Ben Reichardt [†] Email: souravku@usc.edu, Zoom: 8966571313

One-time memory (OTM)	Claim: Octon is i
An ideal one-time memory stores two bits <i>x</i> and <i>y</i> . Bob can choose to read either <i>x</i> or <i>y</i> , but not both.	$\mathbb{P}(\det x) = 1$ \mathbb{F}
$\begin{array}{c} \text{OTM} & \text{get } x \\ x, y & \text{OR} \\ get y \\ \text{NOT BOTH} \end{array}$	OR $\mathbb{P}(\det x) = \frac{3}{4}$
Octon Majorana islands	Store 2 bits in
$X_1 \qquad X_2 $	
Karzig et al. [1] propose that islands of superconducting nanowires correspond to qubits. An "octon" island with 8 Majorana zero modes corresponds to 3 qubits	
 Properties: Only can measure parity of any 2 MZMs The island has overall even parity Two operators commute if they intersect on even number of MZMs (<i>Z</i>₁ and <i>Z</i>₃), 	Alice randomly choose or the X basis (right) to

mperfect OTM

$$\mathbb{P}(\det y) = \frac{3}{4}$$

$$\mathbb{P}(\det y) = 1$$

octon OTM

Read bit from octon OTM

Suppose Bob wants bit *x* (stored in top four MZMs)

He measures parity of bottom four MZMs in either X or Z basis, by measuring the 2 vertical operators or the 2 horizontal operators. If the bottom has even parity, he reads bit x from the horizontal operator. If it has odd parity, he reads bit x from the vertical operator.

Thus, he always obtains bit *x* perfectly. But bit *y* is lost if he chose the wrong basis to get parity of bottom.

$$\mathbb{P}(\det y) = \frac{3}{4}$$



es the Z basis (left) store bits x and y.



A cluster of k octons is equivalent to a nearly perfect OTM. One bit is given by XOR of the top bits of all *k* octons, and the other bit is given by XOR of bottom bits of all *k* octons. If we want to correctly output one bit (say x), the other bit y can be read with probability

$$\mathbb{P}(\operatorname{get} y) = \left(\frac{1}{2} + \frac{1}{2^{k+1}}\right)$$



What about MZM faults?

We choose a CSS code and obtain two classical codes A and B from it.

- In top layer, code A corresponds to the X stabilizers and logical operator X_A .
- In bottom layer, code B corresponds to \bullet Z stabilizers and logical operator Z_B .
- All stabilizer equivalents of X_A and Z_B • intersect.

Thus, we can store the two classical bits as the parity of logical operators X_A and Z_B . Obtaining parity of one logical operator reduces probability of getting the other one.



Conclusion

Best codes when minimum availability of bit x = 0.95



References

- 1. T. Karzig et al., <u>PRB 95, 235305 (2017)</u>
- S. Goldwasser et al., Crypto 2008, 39 2.